2. Ara Değerlendirme - Makale İnceleme Ödevi - Bu ödevin herkes tarafından yapılması Zorunludur

- 2. Ara sınav notunuz sınıfta öncelikle gönüllülük esasına göre atanacak makaleleyi inceleyerek hazırlayacağınız rapor-sunum-kod üzerinden belirlenecek.
- Bu hafta derste olmayanlar için derse geldikleri zaman benim tarafımdan atama yapılacak
- Sunum haftanız makalenin içeriğine bağlı olarak gerekirse kura ile belirlenecek
- En fazla 3 kişiye kadar aynı ödev üzerinde birlikte çalışabilirsiniz - ancak sunum esnasında her birinizin spesifik sorulara cevap vermesi ve herkesin özgün rapor teslim etmesi beklenmektedir.
- Makaleleri okurken yeni terimler ile karşılaşacağınız ve ek referanslara başvurmanız gerekeceğinden şimdiden okumaya başlamanız önemli !!
- Üniversite içindeki bilgisayarlardan/kütüphaneden yararlanmak isteyebileceğiniz pek çok yayına ulaşma olanağınız var
- Hazırladığınız raporlar için
- Ön Teslim tarihi : 21 Nisan – bu tarihe kadar raporun belirli bir bölümü tamamlanmış olmalı
- Son teslim tarihi : 26 Mayıs

2. Ara Değerlendirme - Makale İnceleme Ödevi - Bu ödevin herkes tarafından yapılması Zorunludur

- Ödevlerin önemli bölümü billinen algoritmalardan oluşmaktadır size vereceğim dosyalar içinde örnek kodlar da mevcut bunları kullanabilirsiniz veya internetten ilgili diğer kodlardan yrarlanabilirsiniz. Sunumda sadece kodu çalıştırabiliyor ve anlamış olmanızı bekliyorum.

- Ödevlerle ilgili sınıfta yapacağım ek açıklamaları takip etmek sizin sorumluluğunuzdadır.

- Ödev ve dersle ilgili olarak aşağıdaki linklerden çeşitli araçlara ve matlab kodlarına ulaşabilirsiniz.
  - http://www.cryptool-online.org/index.php?lang=en
  - http://www.cryptool.org/
  - http://www.mathworks.com/matlabcentral

Ödev Listesi – ilk hafta sunulacaklar

00 – Theory and practice of chaotic cryptography – Kaos teorisi temel kavramlar kriptografi alanında uygulamaları

01 – Key distribution and user authentication + Transport-level security – Kitap chapter 4 ve 5 – review kısmındaki noktalar raporda sunuda olmalı *

02 – Wireless network security - Electronic mail security- IP security - Kitap chapter 6,7 ve 8 – review kısmındaki noktalar raporda sunuda olmalı *

03 –System Security - Kitap chapter 9,10 ve 11 – review kısmındaki noktalar raporda sunuda olmalı *

*** Kitap:** Network Securıty Essentıals: *Applıcatıons and Standards* - William Stallings – mail grubunda paylaşıldı

Ödev Listesi – İkinci ve Üçüncü haftalarda sunulacak

04 – IDEA ve PES Blok şifreleme – ana makaleler ve kod verilecek

05 – Camelia ve Blowfish Blok şifreleme – ana makaleler ve kod verilecek

06 – Feal ve Twofish Blok şifreleme – ana makaleler ve kod verilecek

07 – Mars ve Mars II Blok şifreleme – ana makaleler ve kod verilecek

08 – RC5 ve RC6 Blok şifreleme – ana makaleler ve kod verilecek

09 – Serpent ve Skipjack Blok şifreleme – ana makaleler ve kod verilecek

10 – TEA ve Clefia Blok şifreleme – ana makaleler ve kod verilecek

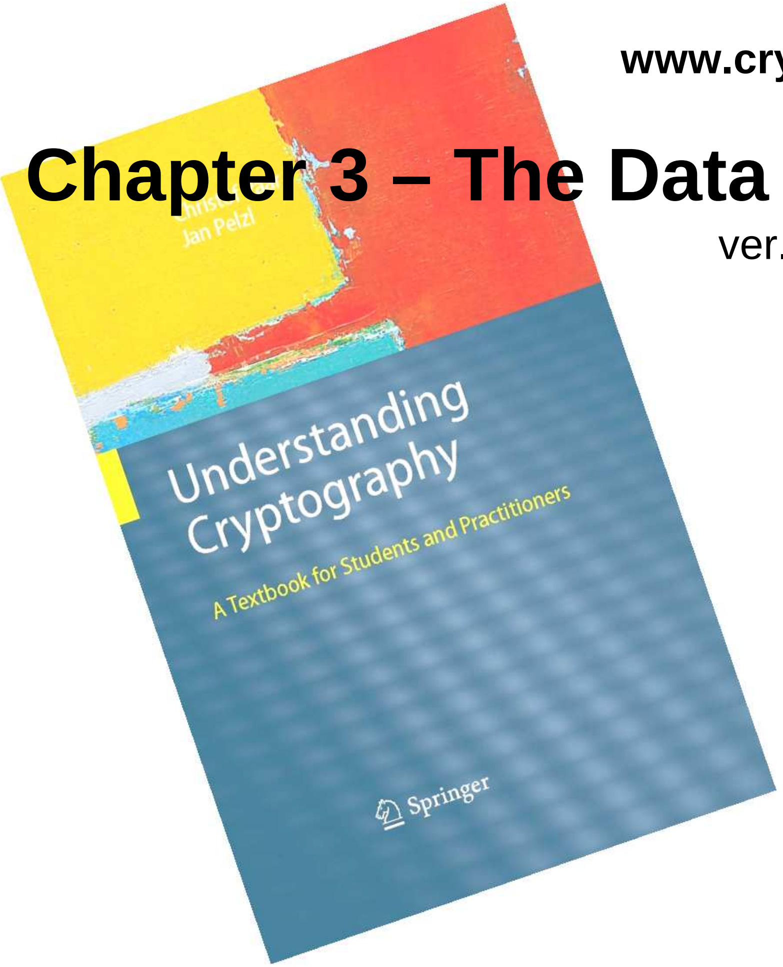
11 – LBlock ve Present Blok şifreleme – ana makaleler ve kod verilecek

12 – Twis ve SEA Blok şifreleme – ana makaleler ve kod verilecek

Understanding Cryptography – A Textbook for Students and Practitioners

by Christof Paar and Jan Pelzl

www.crypto-textbook.com


Christof Paar
Jan Pelzl
Understanding Cryptography
A Textbook for Students and Practitioners
Springer

Chapter 3 – The Data Encryption Standard (DES)

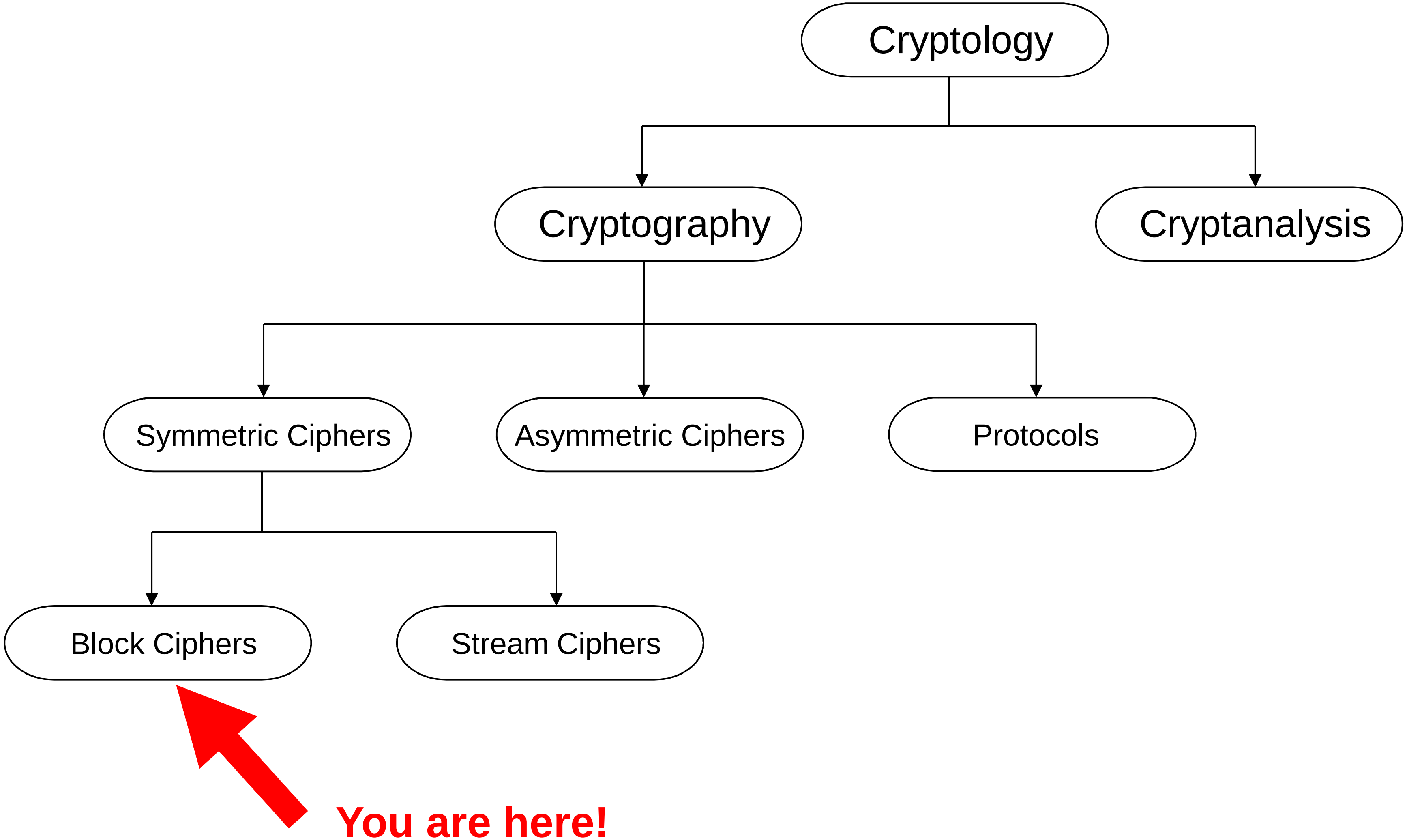
ver. Nov 26, 2010

These slides were prepared by Markus Kasper, Christof Paar and Jan Pelzl

Content of this Chapter

- **Introduction to DES**
- Overview of the DES Algorithm
- Internal Structure of DES
- Decryption
- Security of DES

Classification of DES in the Field of Cryptology



You are here!

DES Facts

- Data Encryption Standard (DES) encrypts **blocks of size 64 bit**.
- Developed by **IBM** based on the cipher *Lucifer* under influence of the *National Security Agency* (NSA), the design criteria for DES have not been published
- **Standardized 1977** by the **National Bureau of Standards** (NBS) today called *National Institute of Standards and Technology* (NIST)
- Most popular **block cipher** for most of the last 30 years.
- By far best studied symmetric algorithm.
- Nowadays considered insecure due to the small **key length of 56 bit.**
- **But: 3DES yields very secure cipher**, still widely used today.
- Replaced by the *Advanced Encryption Standard* (**AES**) in 2000

- For a more detailed history see Chapter 3.1 in *Understanding Cryptography*

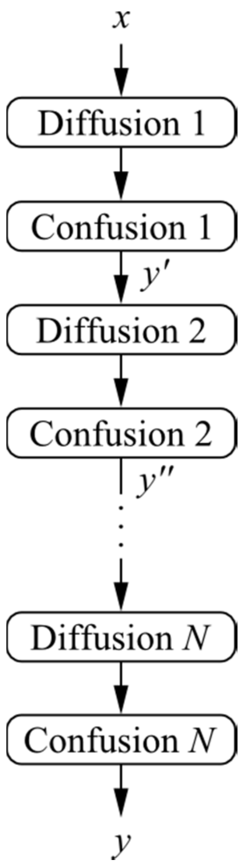
Block Cipher Primitives: Confusion and Diffusion

- Claude Shannon: There are two primitive operations with which strong encryption algorithms can be built:
  1. **Confusion:** An encryption operation where the **relationship between key and ciphertext is obscured**.
     Today, a common element for achieving confusion is **substitution**, which is found in both AES and DES.

  2. **Diffusion:** An encryption operation where the **influence of one plaintext symbol is spread over many ciphertext symbols** with the goal of hiding statistical properties of the plaintext.
     A simple diffusion element is the **bit permutation**, which is frequently used within DES.

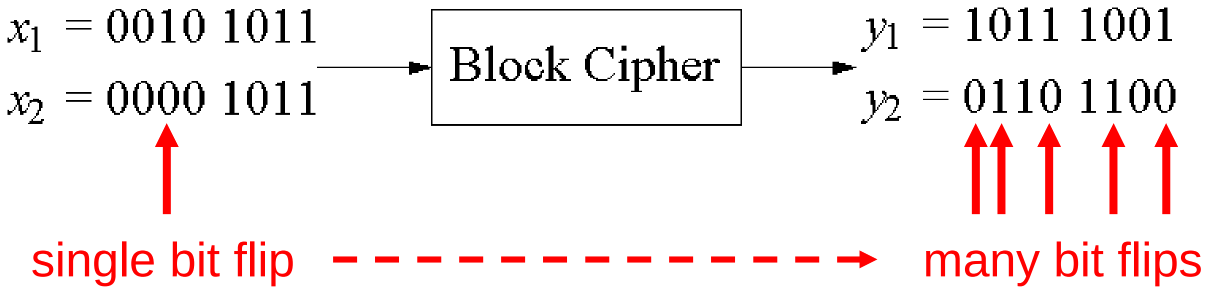
- Both operations by themselves cannot provide security. The idea is to concatenate confusion and diffusion elements to build so called *product ciphers.*

Product Ciphers

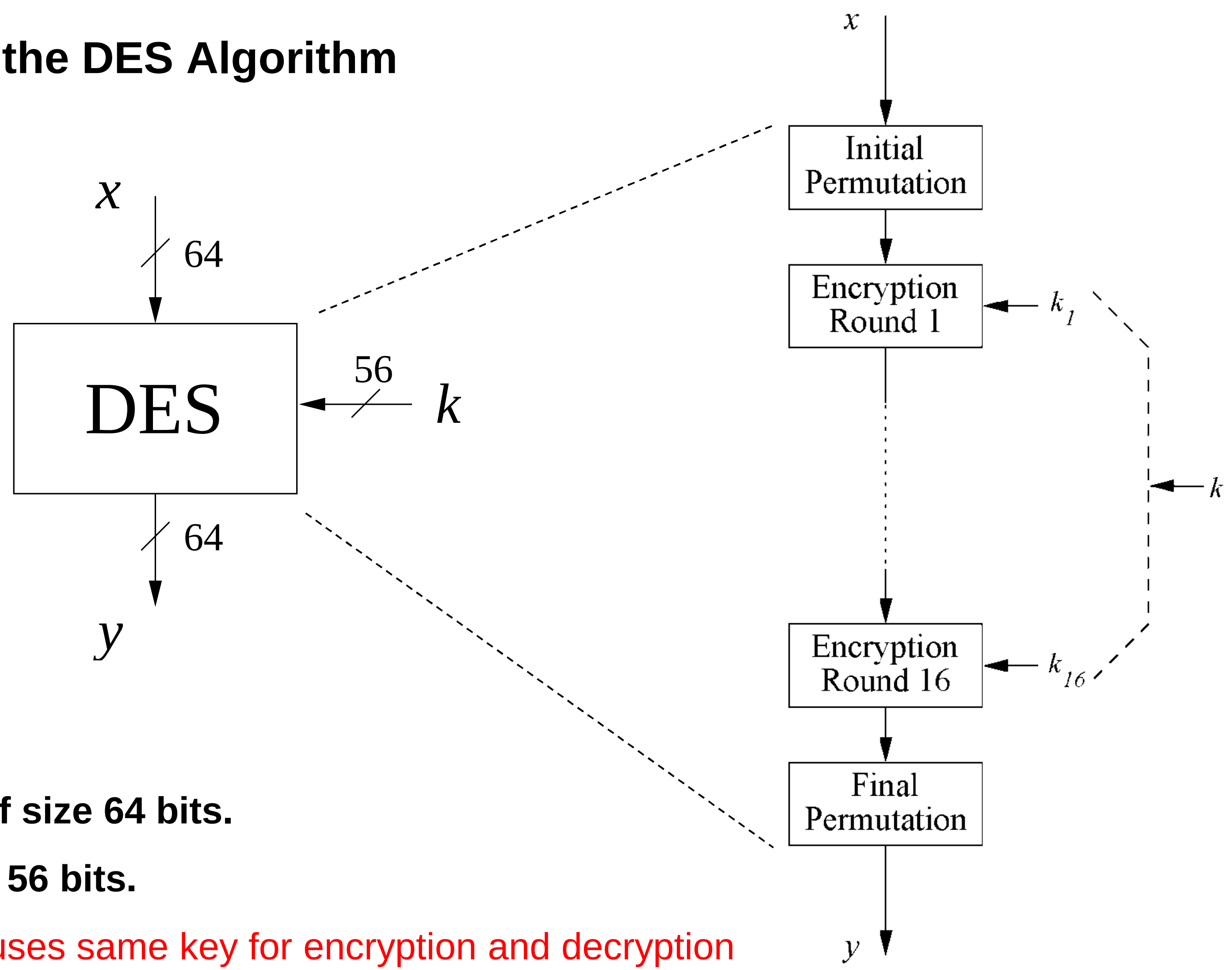


- Most of today‘s block ciphers are *product ciphers* as they consist of rounds which are applied repeatedly to the data.
- Can reach excellent diffusion: **changing of one bit of plaintext results** *on average* in the **change of half the output bits**.

Example:



Overview of the DES Algorithm

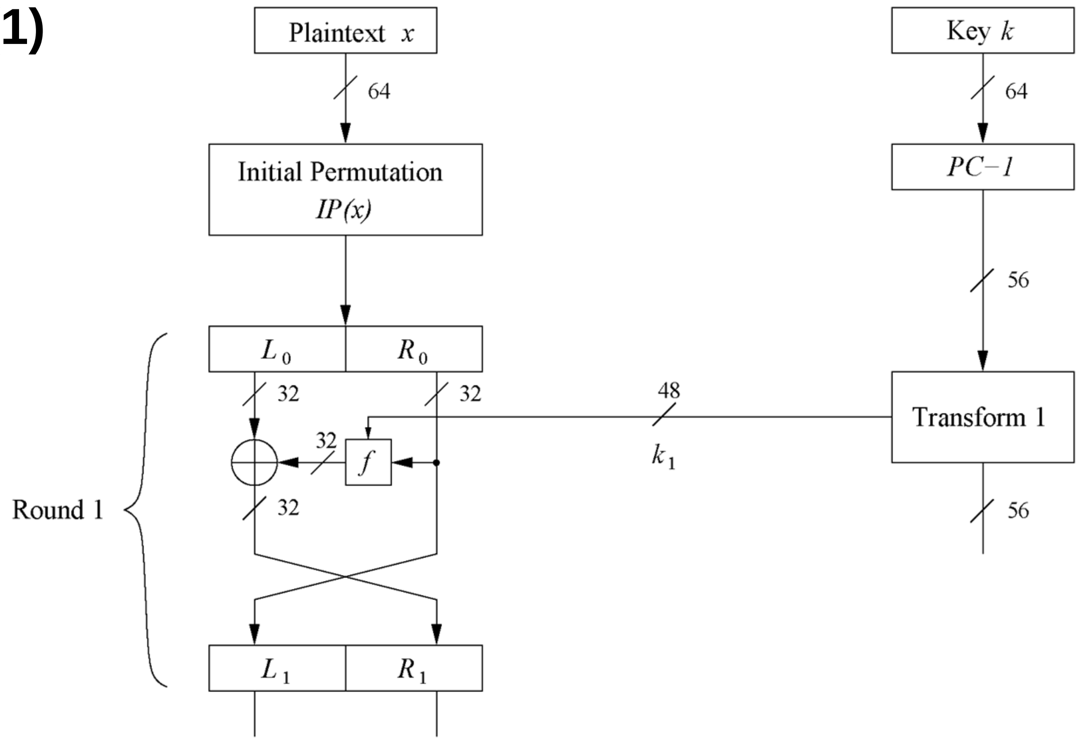


- **Encrypts blocks of size 64 bits.**
- **Uses a key of size 56 bits.**
- Symmetric cipher: uses same key for encryption and decryption
- Uses 16 rounds which all perform the identical operation
- Different subkey in each round derived from main key

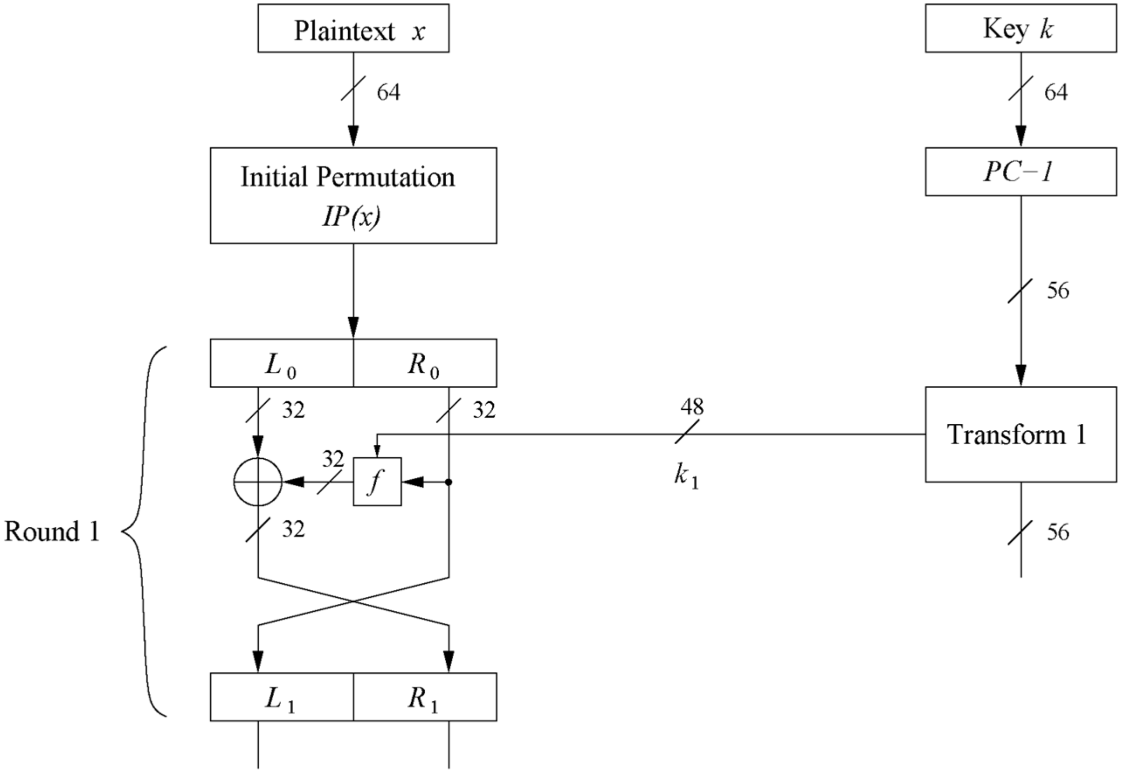
The DES Feistel Network (1)

- DES structure is a *Feistel network*

- Advantage: encryption and decryption differ only in keyschedule
- Deşifreleme sadece tersine dönmüş bir anahtar planı gerektiriyor



- Bitwise initial permutation, then 16 rounds
  1. Plaintext is split into 32-bit halves L_i and R_i
  2. R_i is fed into the function f, the output of which is then XORed with L_i
  3. Left and right half are swapped



Rounds can be expressed as:

$$L_i = R_{i-1},$$
$$R_i = L_{i-1} \oplus f(R_{i-1}, k_i)$$

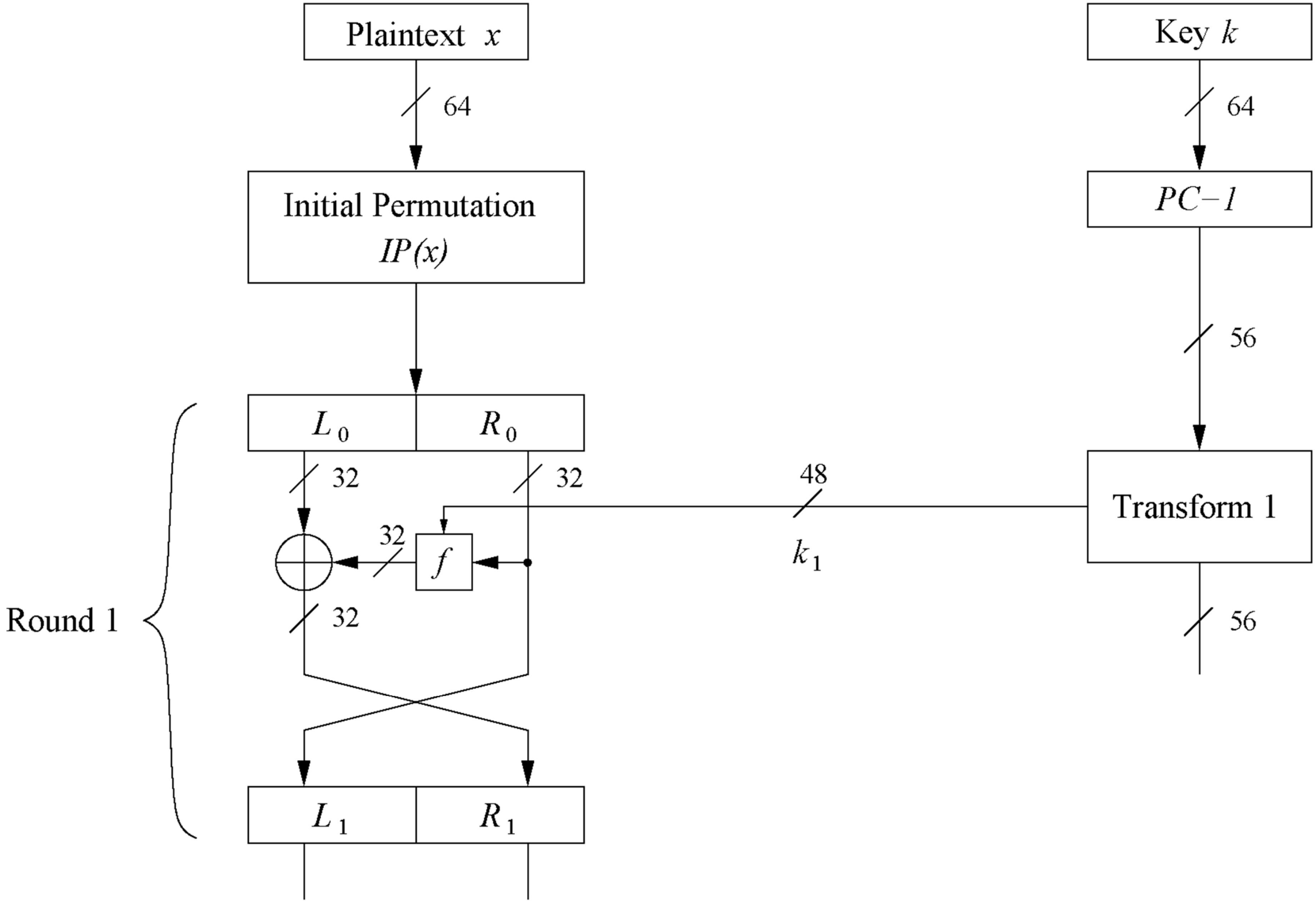
- *IP(x)* initial permutation sonrası şifresiz metin 2 yarıya ayrılıyor ($L_i ve\ R_i$ olarak) bu iki yarı 16 round içeren Fiestel ağına giriyor
- Sağ yarı (R_{i-1}) k_i alt anahtarı ile birlikte f fonksiyonuna girer
- Çıkan 32 bit bu kez sol yarı (L_{i-1}) ile XOR lanır ve bir sonraki roundun sağ yarısını (R_i) oluşturur

$$R_i = L_{i*1} \oplus f(R_{i-1}, k_i)$$

- Sağl yarı hiç değişmeden sonraki round da sağ yarıya aktarılıyor:

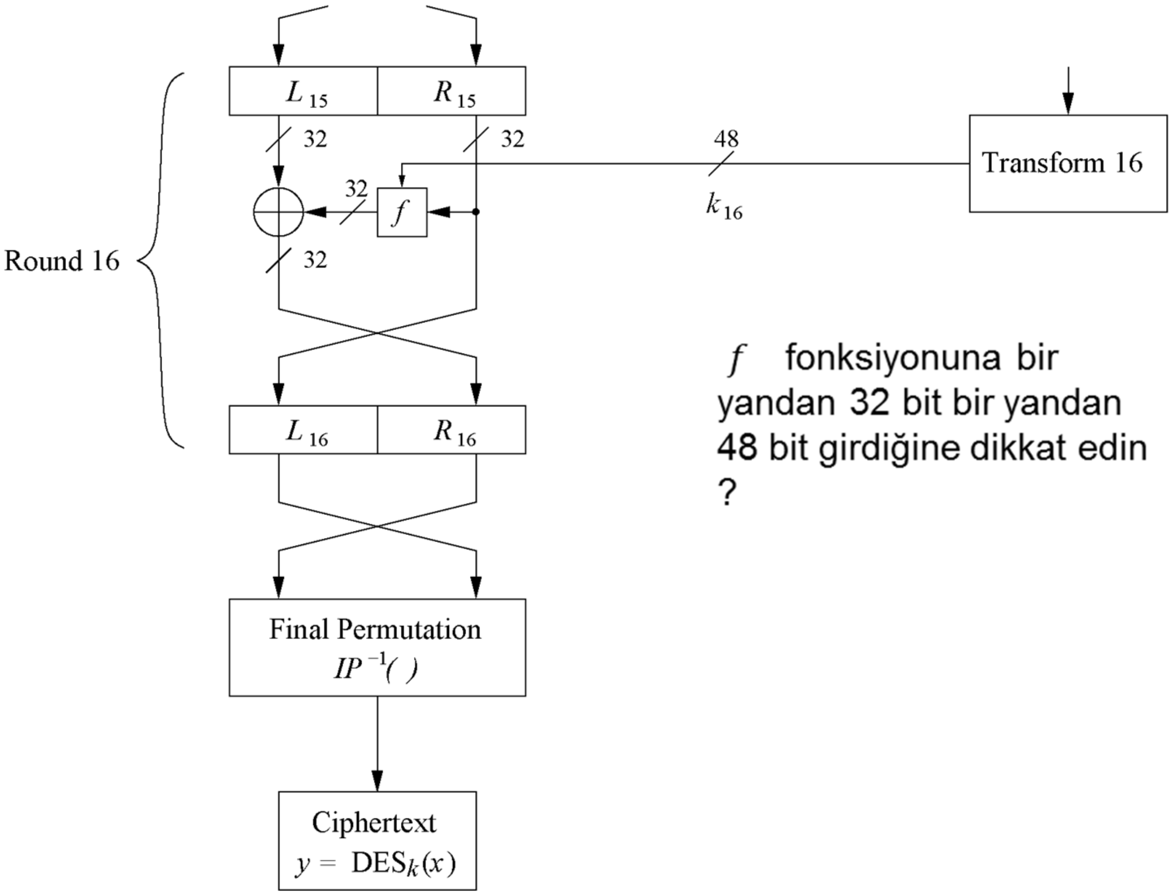
$$L_i = R_{i-1}$$

- PC-1 : Permutated Choice one
- Girişteki 64 bitin 8 bitlik her bloğundan 8. bitleri atarak 56 bite indiriyor
- Dönüşüm işlemi ile üretilen alt anahtar f fonksiyonuna giriyor


Plaintext x
64
Initial Permutation
$IP(x)$
L_0
R_0
32
32
32
f
32
Round 1
L_1
R_1
Key k
64
$PC-1$
56
48
Transform 1
k_1
56

The DES Feistel Network – Son Round

- L and R swapped again at the end of the cipher, i.e., after round 16 followed by a final permutation



16 round sonunda sol ve sağ yarılar $L_{16} e\ R_{16}$ bir kez daha yer değiştirir

Ardından DES in son işlemi olan Final permutation uygulanur

f fonksiyonuna bir yandan 32 bit bir yandan 48 bit girdiğine dikkat edin ?

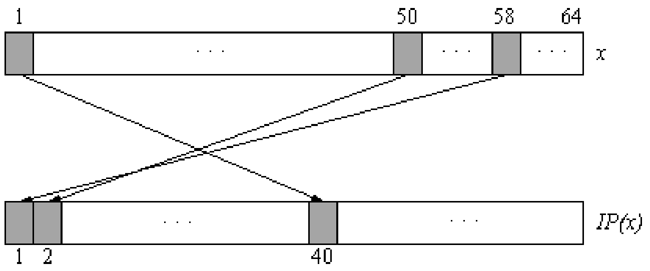
Final permutation gösterimi gibi initial permutation ın tersidir : IP^{-1}

Initial and Final Permutation

- Bitwise Permutations.
- Inverse operations.
- Described by tables *IP* and *IP*$^{-1}$.

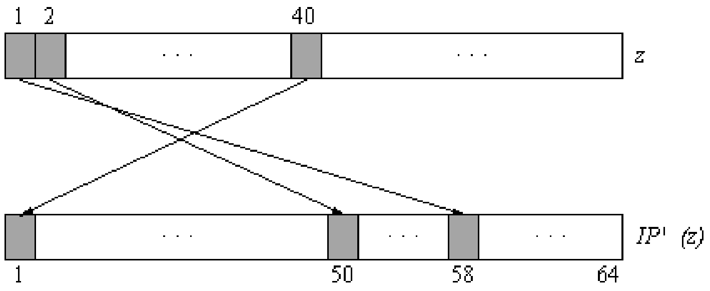
Initial Permutation

IP							
58	50	42	34	26	18	10	2
60	52	44	36	28	20	12	4
62	54	46	38	30	22	14	6
64	56	48	40	32	24	16	8
57	49	41	33	25	17	9	1
59	51	43	35	27	19	11	3
61	53	45	37	29	21	13	5
63	55	47	39	31	23	15	7



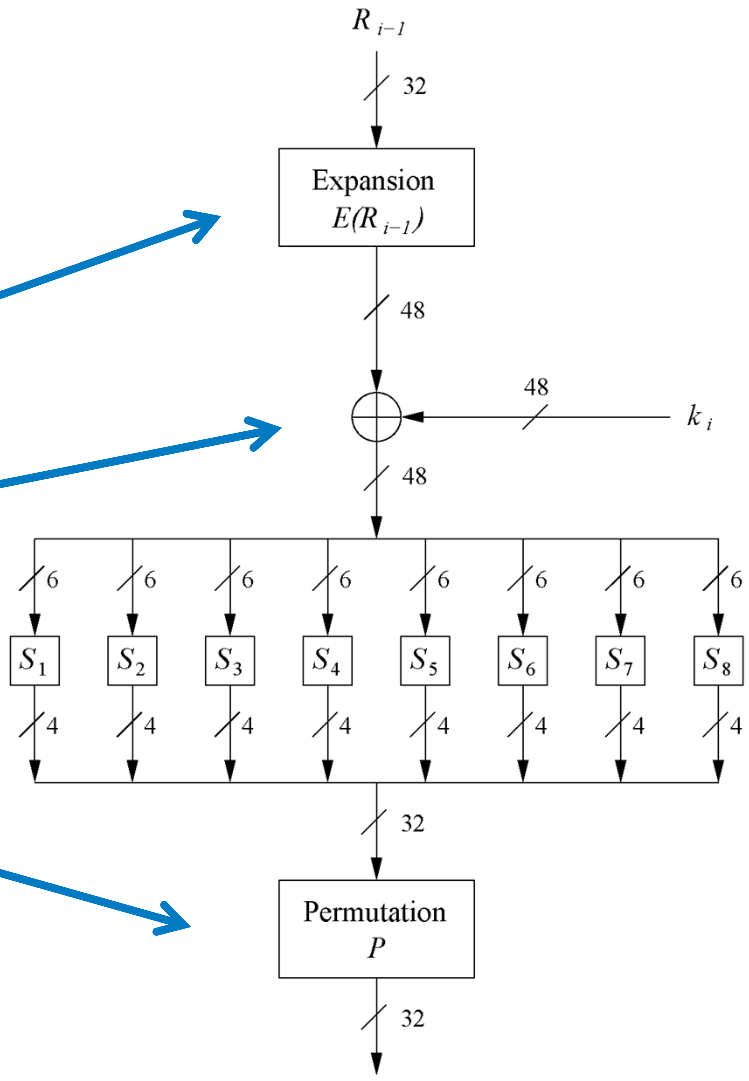
Final Permutation

IP^{-1}							
40	8	48	16	56	24	64	32
39	7	47	15	55	23	63	31
38	6	46	14	54	22	62	30
37	5	45	13	53	21	61	29
36	4	44	12	52	20	60	28
35	3	43	11	51	19	59	27
34	2	42	10	50	18	58	26
33	1	41	9	49	17	57	25



The f-Function

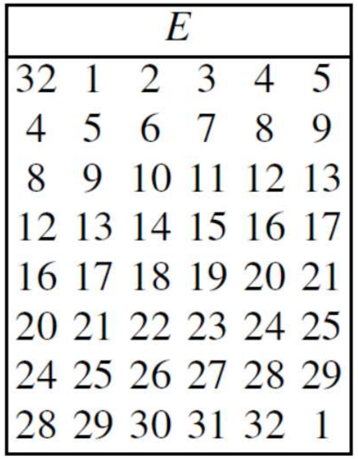
- **main operation of DES**
- *f*-Function inputs:
  R_{i-1} and round key k_i
- **4 Steps**:
  1. Expansion *E*
  2. XOR with round key
  3. S-box substitution
  4. Permutation



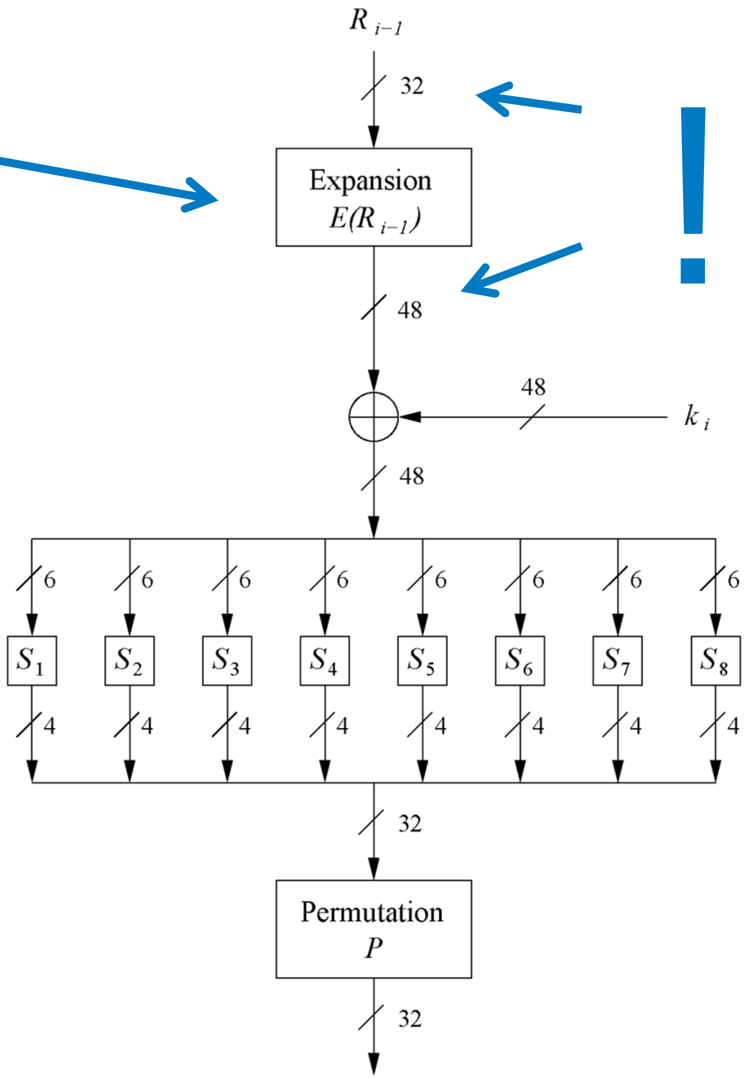
The Expansion Function E

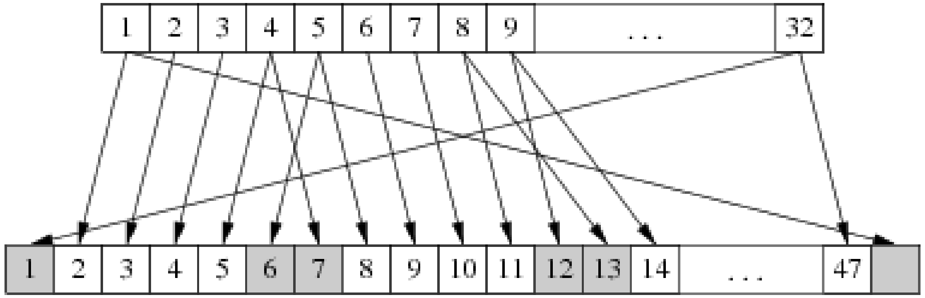
1. **Expansion *E***

- **main purpose:**
  increases diffusion



E					
32	1	2	3	4	5
4	5	6	7	8	9
8	9	10	11	12	13
12	13	14	15	16	17
16	17	18	19	20	21
20	21	22	23	24	25
24	25	26	27	28	29
28	29	30	31	32	1

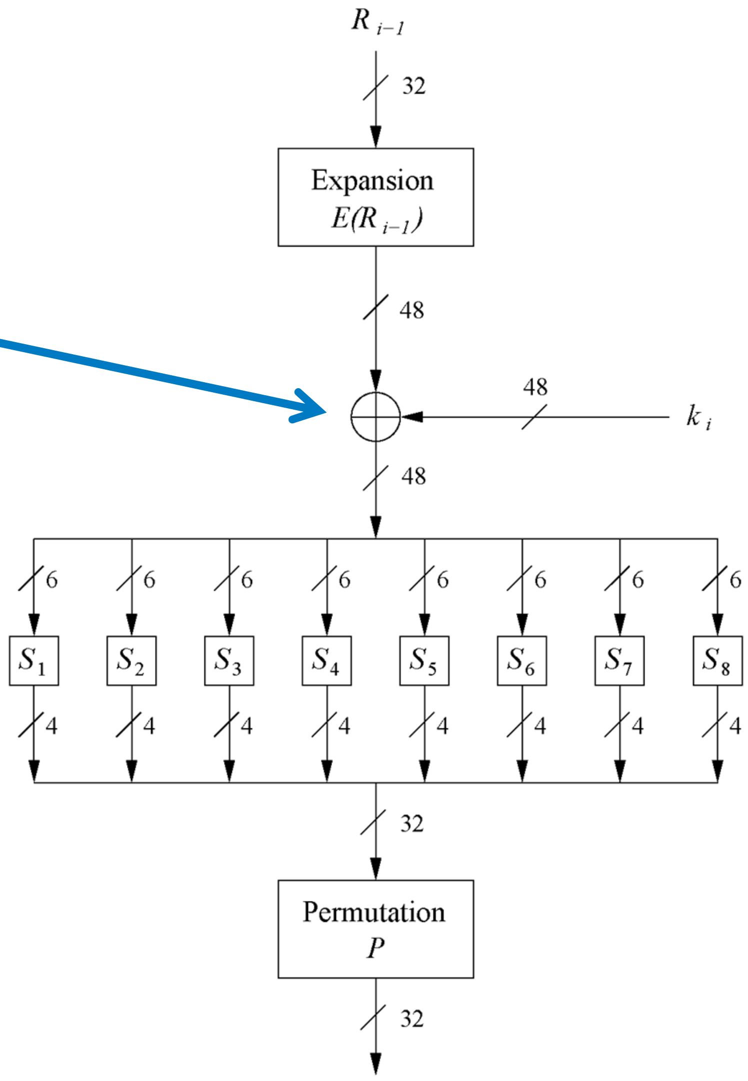




Add Round Key

2. XOR Round Key

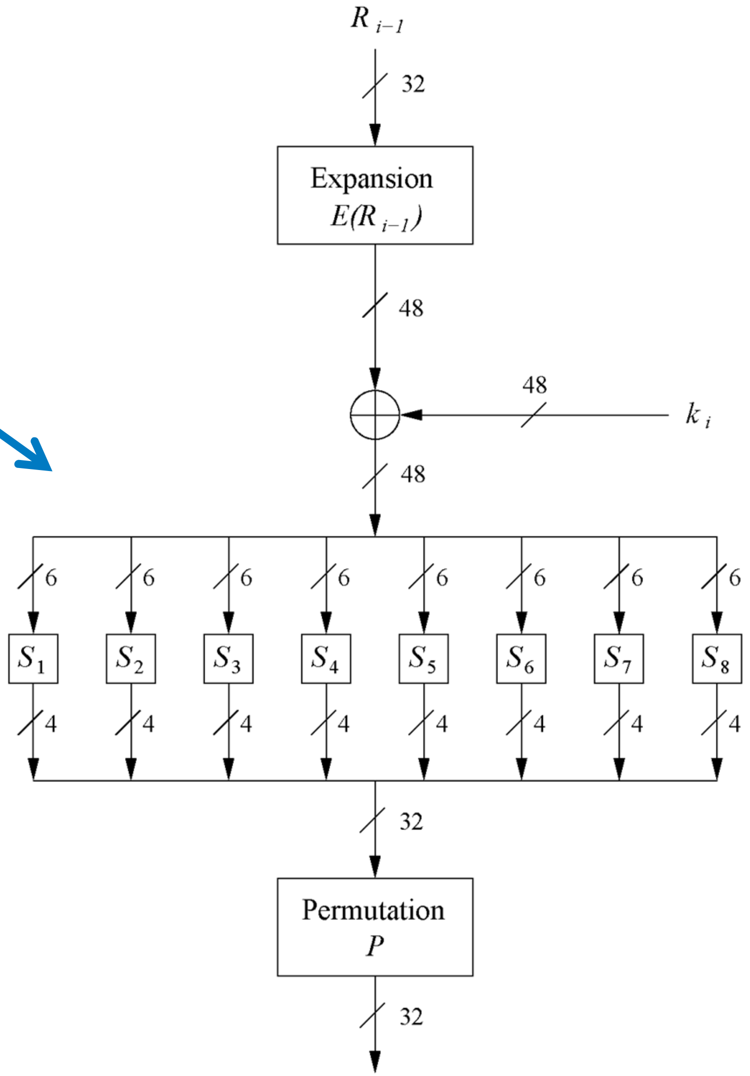
- Bitwise XOR of the round key and the output of the expansion function *E*
- Round keys are derived from the main key in the DES keyschedule (in a few slides)

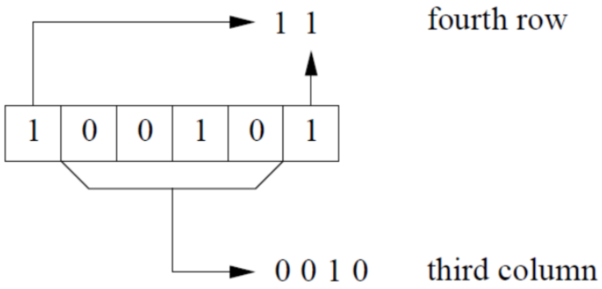


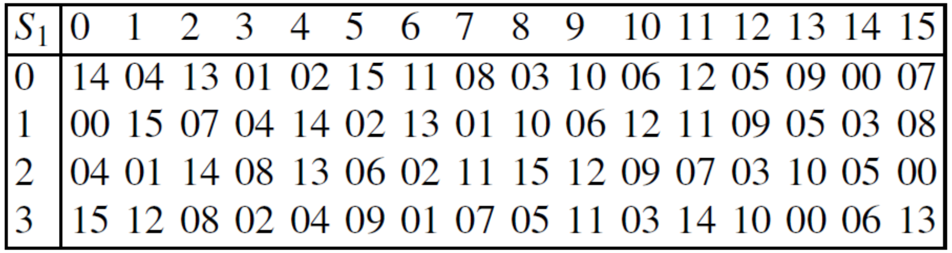
The DES S-Boxes

3. S-Box substitution

- Eight substitution tables.
- 6 bits of input, 4 bits of output.
- Non-linear and resistant to differential cryptanalysis.
- Crucial element for DES security!
- Find all S-Box tables and S-Box design criteria in *Understanding Cryptography* Chapter 3.







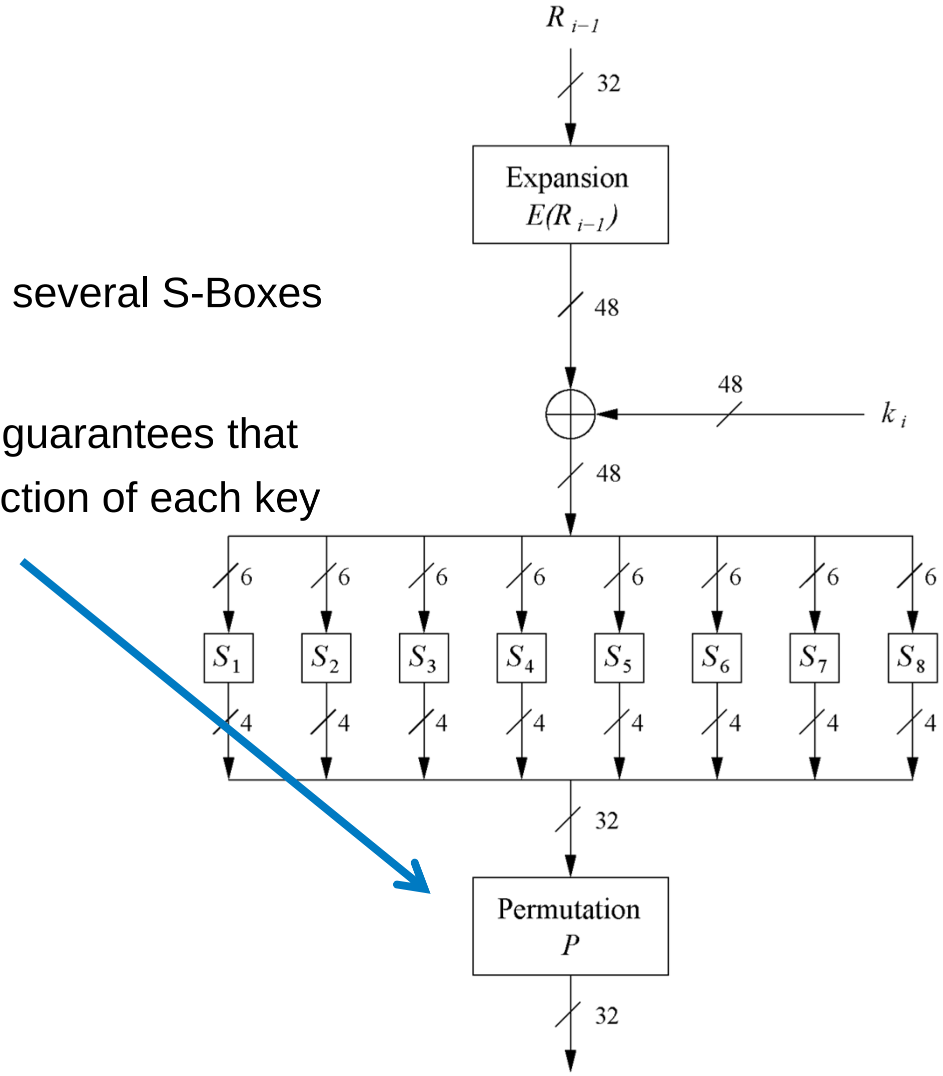
S_1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	14	04	13	01	02	15	11	08	03	10	06	12	05	09	00	07
1	00	15	07	04	14	02	13	01	10	06	12	11	09	05	03	08
2	04	01	14	08	13	06	02	11	15	12	09	07	03	10	05	00
3	15	12	08	02	04	09	01	07	05	11	03	14	10	00	06	13

The Permutation P

4. Permutation P

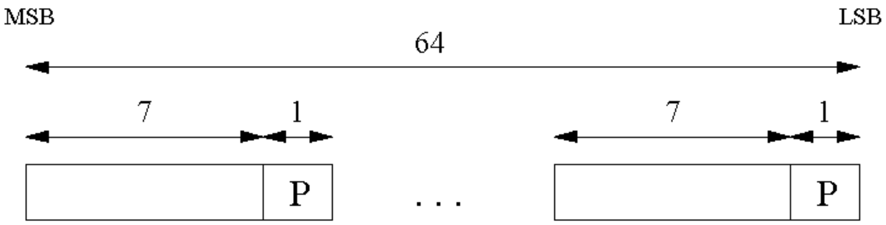
- Bitwise permutation.
- Introduces diffusion.
- Output bits of one S-Box effect several S-Boxes in next round
- Diffusion by E, S-Boxes and P guarantees that after Round 5 every bit is a function of each key bit and each plaintext bit.

P							
16	7	20	21	29	12	28	17
1	15	23	26	5	18	31	10
2	8	24	14	32	27	3	9
19	13	30	6	22	11	4	25



Key Schedule (1)

- Derives 16 round keys (or *subkeys*) k_i of 48 bits each from the original 56 bit key.
- The input key size of the DES is 64 bit: **56 bit key** and 8 bit parity:



P = parity bit

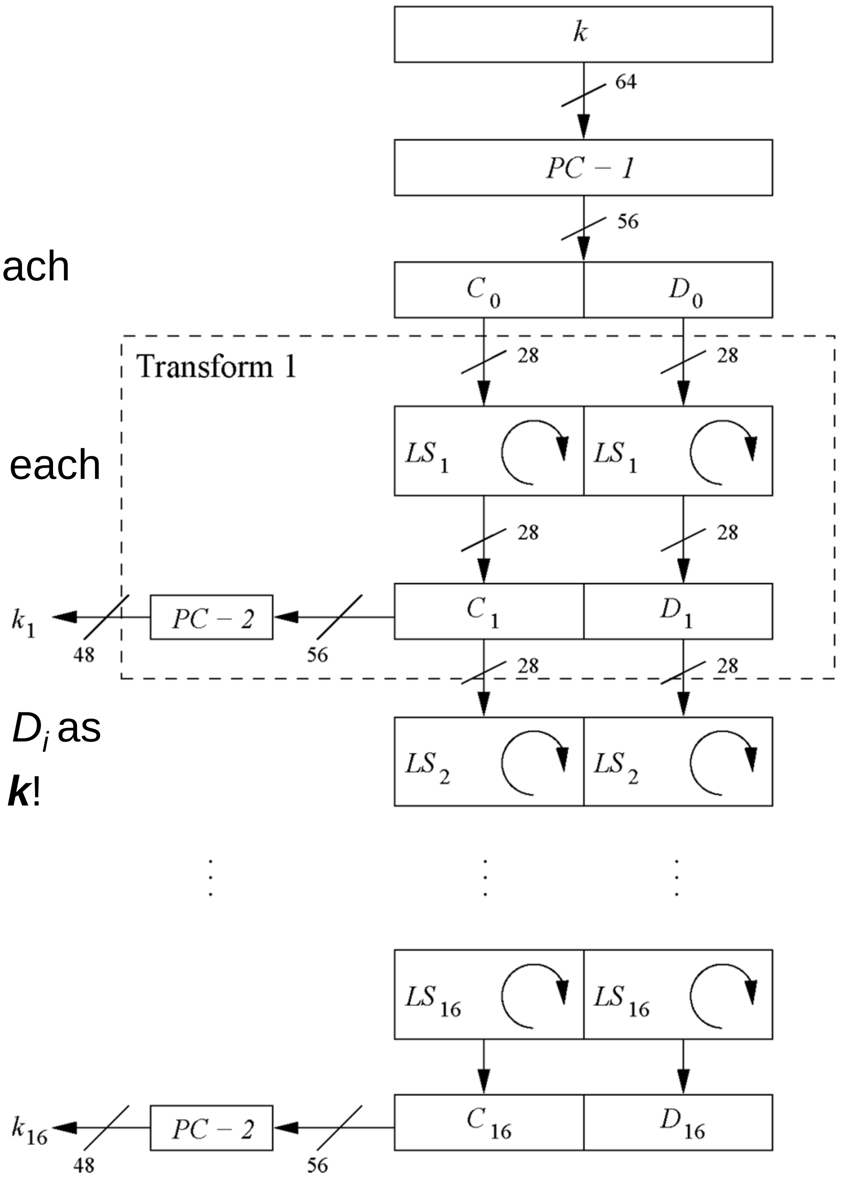
- **Parity bits are removed** in a first **permuted choice** *PC-1*:
  (note that the bits 8, 16, 24, 32, 40, 48, 56 and 64 are not used at all)

$PC-1$							
57	49	41	33	25	17	9	1
58	50	42	34	26	18	10	2
59	51	43	35	27	19	11	3
60	52	44	36	63	55	47	39
31	23	15	7	62	54	46	38
30	22	14	6	61	53	45	37
29	21	13	5	28	20	12	4

Key Schedule (2)

- Split key into 28-bit halves C_0 and D_0.
- In **rounds *i* = 1, 2, 9 ,16,** the two halves are each rotated left by **one bit**.
- In **all other rounds** where the two halves are each rotated left by **two bits**.
- *In each round i permuted choice **PC-2*** selects a permuted subset of 48 bits of C_i and D_i as round key k_i, i.e. **each k_i is a permutation of *k***!

$PC-2$							
14	17	11	24	1	5	3	28
15	6	21	10	23	19	12	4
26	8	16	7	27	20	13	2
41	52	31	37	47	55	30	40
51	45	33	48	44	49	39	56
34	53	46	42	50	36	29	32



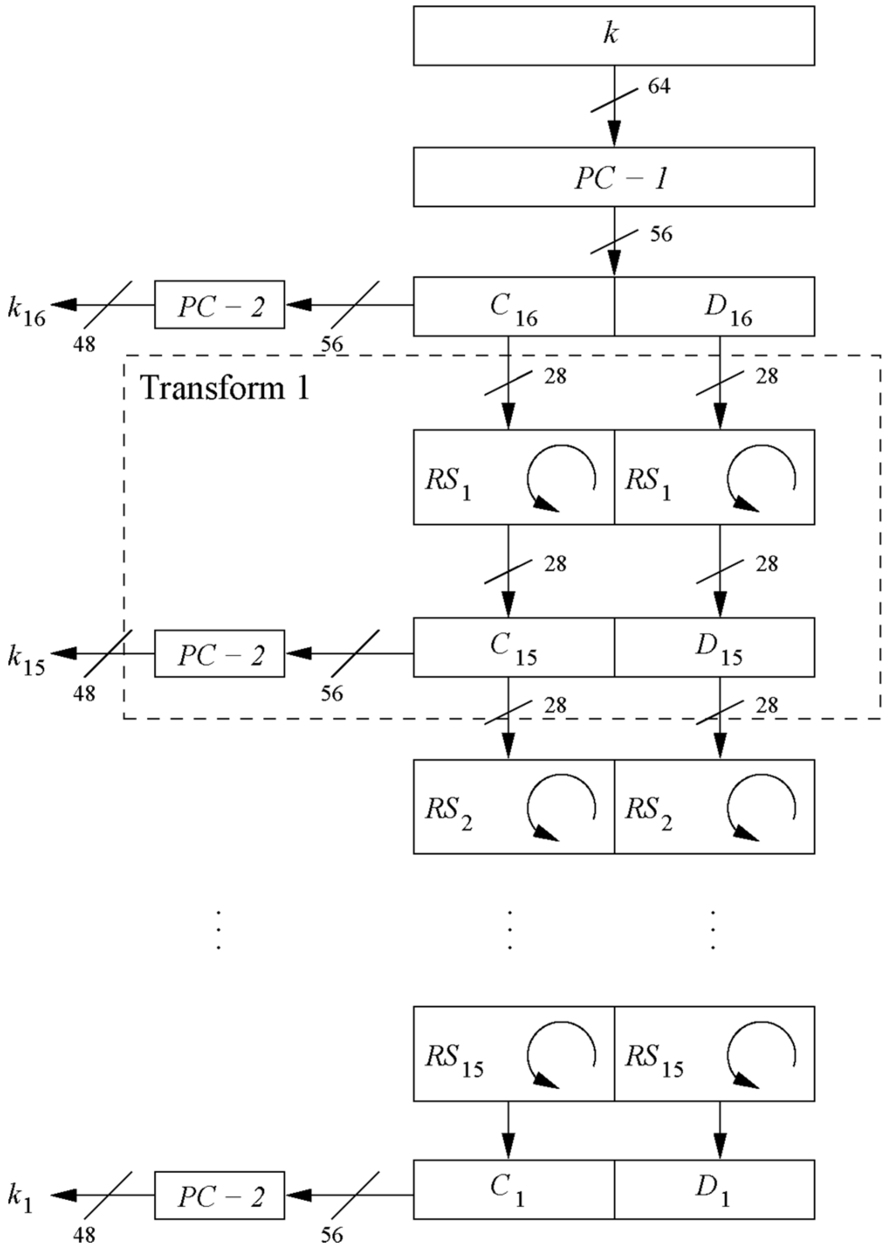
- **Note:** The total number of rotations:
  $4 \times 1 + 12 \times 2 = 28 \Rightarrow D_0 = D_{16}$ and $C_0 = C_{16}$!

Content of this Chapter

- Introduction to DES
- Overview of the DES Algorithm
- Internal Structure of DES
- **Decryption**
- Security of DES

Decryption

- In **Feistel ciphers** only the keyschedule has to be modified for decryption.
- Generate the same 16 round keys in reverse order.
  (for a detailed discussion on why this works see *Understanding Crptography* Chapter 3)
- **Reversed key schedule:**
  As D_0=D_{16} and C_0=C_{16} the first round key can be generated by applying *PC-2* right after *PC-1* (no rotation here!).
  All other rotations of *C* and *D* can be reversed to reproduce the other round keys resulting in:
  - No rotation in round 1.
  - One bit rotation **to the right** in rounds 2, 9 and 16.
  - Two bit rotations **to the right** in all other rounds.



Chapter 3 of *Understanding Cryptography* by Christof Paar and Jan Pelzl

Content of this Chapter

- Introduction to DES
- Overview of the DES Algorithm
- Internal Structure of DES
- Decryption
- **Security of DES**

Security of DES

- **After proposal of DES two major criticisms arose:**
  1. Key space is too small (2^{56} keys)
  2. S-box design criteria have been kept secret: Are there any hidden analytical attacks (*backdoors*), only known to the NSA?

- **Analytical Attacks:** DES is highly resistent to both *differential* and *linear cryptanalysis*, which have been published years later than the DES. This means IBM and NSA had been aware of these attacks for 15 years!
  So far there is no known analytical attack which breaks DES in realistic scenarios.

- **Exhaustive key search:** For a given pair of plaintext-ciphertext (x, y) test all 2^{56} keys until the condition $\text{DES}_k^{-1}(x)=y$ is fulfilled.
  $\Rightarrow$ Relatively easy given today's computer technology!

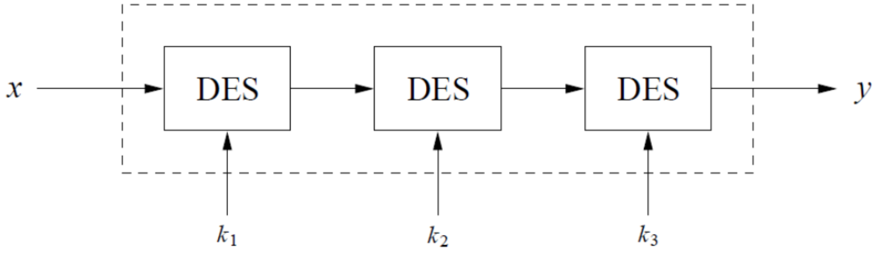
History of Attacks on DES

Year	Proposed/ implemented DES Attack
1977	Diffie & Hellman, (under-)estimate the costs of a key search machine
1990	Biham & Shamir propose differential cryptanalysis (2^{47} chosen ciphertexts)
1993	Mike Wiener proposes design of a very efficient key search machine: Average search requires 36h. Costs: $1.000.000
1993	Matsui proposes linear cryptanalysis (2^{43} chosen ciphertexts)
Jun. 1997	DES Challenge I broken, 4.5 months of distributed search
Feb. 1998	DES Challenge II--1 broken, 39 days (distributed search)
Jul. 1998	DES Challenge II--2 broken, key search machine *Deep Crack* built by the Electronic Frontier Foundation (EFF): 1800 ASICs with 24 search engines each, Costs: $250 000, 15 days average search time (required 56h for the Challenge)
Jan. 1999	DES Challenge III broken in 22h 15min (distributed search assisted by *Deep Crack*)
2006-2008	Reconfigurable key search machine *COPACOBANA* developed at the Universities in Bochum and Kiel (Germany), uses 120 FPGAs to break DES in 6.4 days (avg.) at a cost of $10 000.

Triple DES – 3DES

- Triple encryption using DES is often used in practice to extend the effective key length of DES to 112. For more info on multiple encryption and effective key lengths see Chapter 5 of *Understanding Cryptography.*

$$y = DES_{k_3}(DES_{k_2}(DES_{k_1}(x)))$$



- Alternative version of *3DES:* $y = DES_{k_3}(DES^{-1}_{k_2}(DES_{k_1}(x))).$

  Advantage: choosing k_1=k_2=k_3 performs single DES encryption.
- No practical attack known today.
- Used in many legacy applications, i.e., in banking systems.

Alternatives to DES

Algorithm	I/O Bit	key lengths	remarks
AES / Rijndael	128	128/192/256	DES "replacement", worldwide used standard
Triple DES	64	112 (effective)	conservative choice
Mars	128	128/192/256	AES finalist
RC6	128	128/192/256	AES finalist
Serpent	128	128/192/256	AES finalist
Twofish	128	128/192/256	AES finalist
IDEA	64	128	patented

Lessons Learned

- DES was the dominant symmetric encryption algorithm from the mid-1970s to the mid-1990s. Since 56-bit keys are no longer secure, the Advanced Encryption Standard (AES) was created.
- Standard DES with 56-bit key length can be broken relatively easily nowadays through an exhaustive key search.
- DES is quite robust against known analytical attacks: In practice it is very difficult to break the cipher with differential or linear cryptanalysis.
- By encrypting with DES three times in a row, triple DES (3DES) is created, against which no practical attack is currently known.
- The “default” symmetric cipher is nowadays often AES. In addition, the other four AES finalist ciphers all seem very secure and efficient.